

Application of Sawi transform in Cryptography

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Abstract

Cryptography is the science of secret communication. Encryption is a process used to convert the given original text into a secret text known as cipher text. Decryption is a process of converting the cipher text into the original text. In this paper, we discuss an encryption and decryption algorithm based on Sawi transform. We illustrate the algorithms by an example.

Keywords: Sawi transform, encryption, decryption, plain text, cipher text.

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Introduction

Cryptography is the science of secret and secured communication [1]. Encryption is a process of encoding (converting) a given original text (called plain text) into a cipher text. Decryption is a process of decoding (re-converting) the cipher text into the original text [4]. Various cryptographic techniques based on mathematical algorithms have been introduced [1, 4]. Cryptography is used various fields (including agriculture) where data security is important.

Integral transforms has its applications in various fields. Recently, integral transforms has been employed in cryptography [2, 3]. Sawi transform transform was introduced by Mohand Mahgoub [5]. In this paper, we provide an encryption and a decryption algorithm based on Sawi transform and affine cipher. We also illustrate the proposed algorithm by an example.

Sawi Transform

In this section, we give the definition and basic properties of Sawi transform [5].

Let

$$A = \{f(t) : \exists M, k_1, k_2 > 0 \ni |f(t)| < M e^{\frac{|t|}{k_j}}, \text{ for } t \in (-1)^j \times [0, \infty)\}$$

where M is a finite constant and k_1, k_2 may be finite or infinite. For $f(t) \in A$, the Sawi transform $S(\cdot)$ is defined by

$$S[f(t)] = R(v) = \frac{1}{v^2} \int_0^{\infty} f(t) e^{-t/v} dt, \quad k_1 \leq v \leq k_2$$

provided the integral exists.

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Remark: The sufficient conditions for the existence of Sawi transform are that $f(t)$ for $t \geq 0$ be piecewise continuous and of exponential order.

Sawi transform of some basis functions

1. $f(t) = 1, S[f(t)] = R(v) = \frac{1}{v}.$
2. $f(t) = t, S[f(t)] = R(v) = 1.$
3. For $n \geq 2, f(t) = t^n, S[f(t)] = R(v) = v^{n-1}n!.$

Inverse Sawi transform

If $S[f(t)] = R(v)$ then the inverse sawi transform is $S^{-1}(R(v)) = f(t).$

$$S^{-1}\left[\frac{1}{v}\right] = 1. \quad 1.$$

$$S^{-1}[1] = t. \quad 2.$$

$$3. S^{-1}[v^{n-1}] = t^n/n! \text{ for } n \geq 1.$$

Property

Sawi transform satisfies the linearity property.

$$S[af(t) + bg(t)] = aS[f(t)] + bS[g(t)].$$

Main Results

In this section, we give an encryption and a decryption algorithm based on sawi transform and affine cipher.

Using the table below, we convert the alphabets into numbers and viceversa.

A	B	C	D	E	F	G	H	I	J	K	L	M
1	2	3	4	5	6	7	8	9	10	11	12	13
N	O	P	Q	R	S	T	U	V	W	X	Y	Z
14	15	16	17	18	19	20	21	22	23	24	25	26

Here upper case and lower case alphabets are undifferentiated, special characters and spaces are not considered.

Encryption Algorithm

- (i) Assign every alphabet in the plain text to a number.
- (ii) Corresponding to the plain text consider the

$$E(x) = ax + b \pmod{26},$$

where a and b are keys for ciphers and a is co-prime to 26.

- (iii) Let $p(t)$ be the polynomial of degree $n - 1$ with coefficients as $E(x)$. Here n is the number of terms in the plain text.

- (iv) Apply Sawi transform to $p(t)$.

$$n-1$$

- (v) Let $S[p(t)] = R(v) = \sum_{i=0}^n q_i v^{i-1}.$

(vi) Write $q_i = 26c_i + r_i$ for each i with $0 < r_i < 26$.
 Here c_i 's are the keys.

(vii) By converting each r_i to its corresponding alphabet, we can get the cipher text.

Decryption Algorithm

Assume the secret keys a, b are known to the receiver.

Consider the cipher text with the keys c_0, c_1, \dots, c_{n-1} .

1. Convert the cipher text into the number sequence r_0, r_1, \dots, r_{n-1} .
2. Find $q_i = 26c_i + r_i$ for each $i = 0, 1, 2, \dots, n - 1$.

$$n-1$$
3. Let $R(v) = \sum_{i=0}^{n-1} q_i v^{i-1}$.
4. Apply inverse Sawi transform to $R(v)$ and get the polynomial $p(t)$ of degree $n - 1$.
5. Consider the coefficients of $p(t)$ as a number sequence. 6. For each number y in the number sequence, calculate

$$D(y) = a^{-1}(y - b) \pmod{26}.$$

Here a^{-1} denotes the multiplicative inverse of 'a' under modulo 26.

7. By converting the resulting number sequence into alphabets, we can get the original text.

Illustrative Example

Consider the plain text

“DELIGHT”

Take $a = 11$ and $b = 13$. Then clearly 11 is co-prime to 26 and $a^{-1} = 19$.

Encryption

1. The number sequence corresponding the plain text after applying affine cipher is given below.

Plain text	D	E	L	I	G	H	T
Number sequence	4	5	12	9	7	8	20
$ax + b = 11x + 13$	57	68	145	112	90	101	233
$E(x) = 11x + 13 \pmod{26}$	5	16	15	8	12	23	25

Therefore, the number sequence is 5, 16, 15, 8, 12, 23, 25.

2. Here $n = 7$.

Let $p(t) = 5 + 16t + 15t^2 + 8t^3 + 12t^4 + 23t^5 + 25t^6$.

3. By applying Sawi transform to $p(t)$, we get

$$\begin{aligned}
 S[p(t)] = R(v) &= S[5 + 16t + 15t^2 + 8t^3 + 12t^4 + 23t^5 + 25t^6] \\
 &= S[5] + 16S[t] + 15S[t^2] + 8S[t^3] + 12S[t^4] + 23S[t^5] + 25S[t^6] \\
 &= 5 \frac{1}{v} + 16 + 15 \times 2!v + 8 \times 3!v^2 + 12 \times 4!v^3 + 23 \times 5!v^4 + 25 \times 6!v^5 \\
 &= \frac{5}{v} + 16 + 30v + 48v^2 + 288v^3 + 2760v^4 + 18000v^5
 \end{aligned}$$

Thus $q_0 = 5, q_1 = 16, q_2 = 30, q_3 = 48, q_4 = 288, q_5 = 2760, q_6 = 18000$

4. For each i , write $q_i = 26c_i + r_i$. We get $c_0 = 0, c_1 = 0, c_2 = 1, c_3 = 1, c_4 = 11, c_5 = 106, c_6 = 692$ and $r_0 = 5, r_1 = 16, r_2 = 4, r_3 = 22, r_4 = 2, r_5 = 4, r_6 = 8$

5. Convert the number sequence $r_0, r_1, r_2, r_3, r_4, r_5, r_6$ into alphabets

5	16	4	22	2	4	8
E	P	D	V	B	D	H

Thus the cipher text is "EPDVBDH"

Decryption

Consider the cipher text "EPDVBDH" with the keys $c_0 = 0, c_1 = 0, c_2 = 1, c_3 = 1, c_4 = 11, c_5 = 106, c_6 = 692$

1. The number sequence corresponding to the cipher text is $r_0 = 5, r_1 = 16, r_2 = 4, r_3 = 22, r_4 = 2, r_5 = 4, r_6 = 8$
2. For each i , write $q_i = 26c_i + r_i$. Then we get $q_0 = 5, q_1 = 16, q_2 = 30, q_3 = 48, q_4 = 288, q_5 = 2760, q_6 = 18000$

3. Take $R(v) = \sum_{i=0}^n q_i v^{i-1}$. Then

$$R(v) = \frac{5}{v} + 16 + 30v + 48v^2 + 288v^3 + 2760v^4 + 18000v^5$$

4. Apply inverse Sawi transform to $R(v)$, we get

$$S^{-1}[R(v)] = p(t) = 5 + 16t + 15t^2 + 8t^3 + 12t^4 + 23t^5 + 25t^6.$$

5. Consider the coefficients as a number sequence. Then we get the number sequence as 5, 16, 15, 8, 12, 23, 25.

6. For each y in the number sequence, calculate $D(y) = 19(y-13) \pmod{26}$

	5	16	15	8	12	23	25
$19(y - 13)$	-152	57	38	-95	-19	190	228
$D(y) = 19(y - 13) \pmod{26}$	4	5	12	9	7	8	20
Plain text	D	E	L	I	G	H	T

Thus the plain text is "DELIGHT"

Conclusion

In cryptography, mathematical techniques have been employed to secure messages. In this paper, we have introduced an encryption algorithm using Sawi transform and a decryption algorithm using inverse Sawi transform.. We illustrated the proposed algorithm by an example.

References

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